

Scientific Computing on Parallel Machines

Cálculo científico en ordenadores paralelos

Markus Uhlmann

CIEMAT

www.ciemat.es/sweb/comfos/personal/uhlmann

UCM – Abril 2008

Schedule

Part I	Introduction to parallel programming	lecture 1
Part II	Introduction to MPI	
	General introduction	
	& MPI program structure	lecture 2
	Point-to-point communication	lecture 3,4
	Collective communication	lecture 5
	2D Poisson example	lecture 6,7
	Non-contiguous data & Mixed datatypes	lecture 7
	Virtual topologies & Communication subsets	lecture 8
	Use of linear algebra libraries	lecture 9
	Extensions	lecture 10

Case Study: Jacobi Iteration of 2D Poisson Problem

$$(\partial_{xx} + \partial_{yy}) u = 0$$

$$u(\mathbf{x} \in \Gamma) = u_\Gamma$$

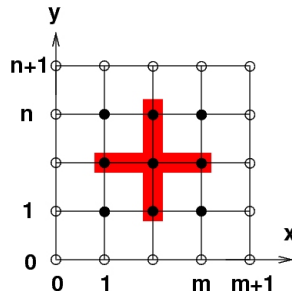
here: $u(x, y) = \sin(\pi x) \cdot \exp(\pi y)$

- 2nd order central finite-differences, uniform grid:

$$\frac{u_{i+1,j} + u_{i-1,j} + u_{i,j+1} + u_{i,j-1} - 4u_{i,j}}{h^2} = 0$$

Jacobi iteration method:

$$u_{i,j}^{n+1} = \frac{u_{i+1,j}^n + u_{i-1,j}^n + u_{i,j+1}^n + u_{i,j-1}^n}{4}$$



Jacobi Iteration of 2D Poisson Problem

Sequential code structure

- ▶ initialisation
 - ▶ iterative loop
 - ▶ advance the solution by one step
 - ▶ test for convergence
 - ▶ compute error w.r.t. analytical solution
 - ▶ write solution to file
- ⇒ useful to employ modular structure
(sequential source code)

Parallel Code

Individual steps

1. identify concurrency
2. decomposition: (domain/functional)
3. design of communication patterns
4. algorithm implementation
5. validation against sequential results
6. timing/optimization of parallel execution

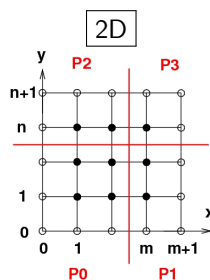
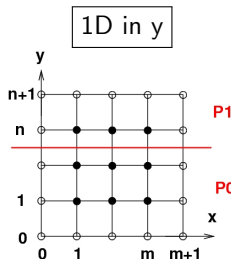
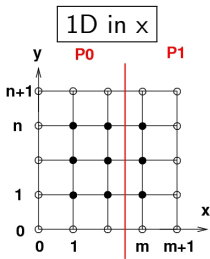
Concurrency/communication

Symbolic code

```
(1) initialisation  
for each iteration  
    (2a) update the solution  
    (2b) check for convergence  
end
```

- ▶ main work load
 - ▶ potential for concurrency
- ⇒ domain decomposition
- starting point: main program template with utilities

Cartesian Domain Decomposition



messages: $(np - 1)^2$

$(np - 1)^2$

size: $(np - 1)^2 n$

$(np - 1)^2 m$

$(n_{xp} - 1)^2 n_{yp}$
 $+ (n_{yp} - 1)^2 n_{xp}$
 $(n_{xp} - 1)^2 n$
 $+ (n_{yp} - 1)^2 m$

- for large np : 2D decomposition $\Rightarrow$ more messages, smaller total volume

Domain/Data Decomposition

- ▶ 2D decomposition is most flexible
- ▶ BUT: for simplicity we choose 1D decomposition
- ▶ no decomposition for first array index
→ contiguous data (FORTRAN)

1D Decomposition in Y-Direction: Pointers

Relation between global and local indices

⇒ subroutine defining pointer arrays (known to all processes):

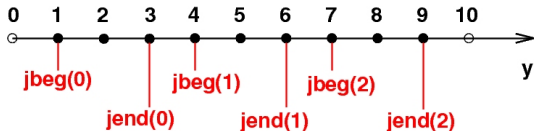
```
integer jbeg(0:nproc-1)  
integer jend(0:nproc-1)
```

example:

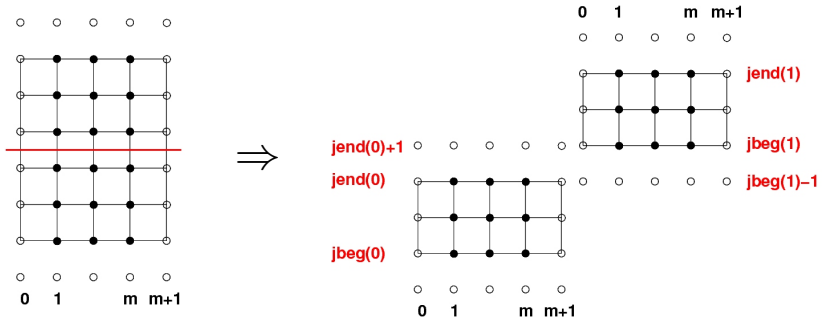
▶ $nproc=3$
 $n=9$

▶ pointer routine: code template

→ possible solution



Ghost Cells



⇒ arrays have size:
equivalently:

$A(0:m+1, jbeg(myrank)-1:jend(myrank)+1)$
 $A(0:m+1, 1:jend(myrank)-jbeg(myrank)+1)$

Communication

Symbolic code

```
(1) initialisation
for each iteration
  (2a) update the solution
  (2b) check for convergence
toto
end
```

```
(1) initialisation
for each iteration
  (2a) update local part of solution
  (2b) check for convergence
  (2c) exchange ghost-cell data
end
```

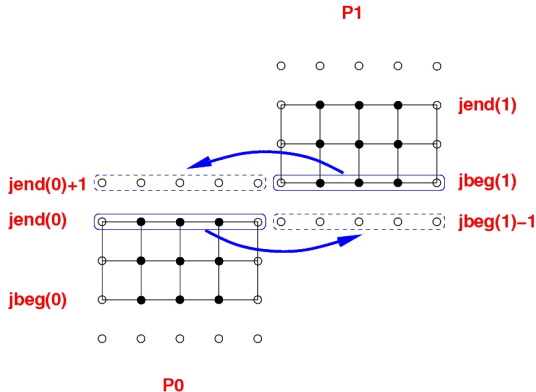
Convergence criterion: communication

Compute global maximum of increment

Collective communication

- ▶ reduction operation
 - ▶ `MPI_MAX`
 - ▶ all-to-all (all tasks need the result)
- ⇒ `MPI_ALLREDUCE`
- possible solution
- ~> how could this routine be improved? (latency)

Ghost Cells: Data Exchange



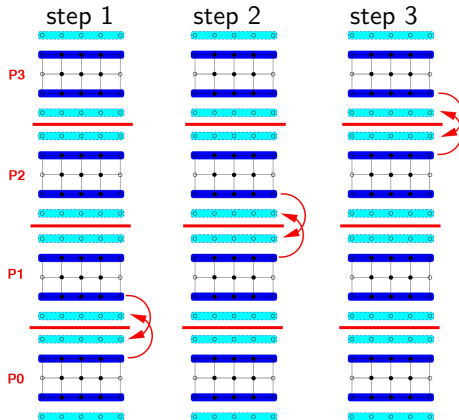
- code template for communication routine

Ghost Cells: Data Exchange

Communication patterns depend on several choices

- ▶ blocking vs. non-blocking
- ▶ send & recv vs. sendrecv
- ▶ order in which communication ops are posted

Communication Pattern 1



- blocking communication
- inferior/superior order
→ 'domino' effect:
⇒ needs $(nproc-1)$ steps

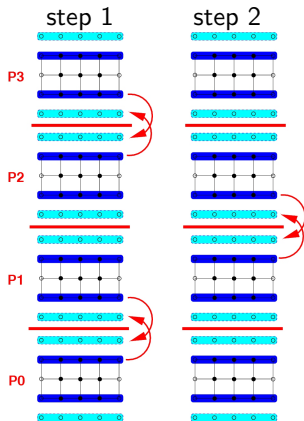
(source code – Fortran)

Communication Pattern 2

non-blocking send/receive:

- ▶ each processor posts 2 send and 2 receive operations
- ▶ MPI internally decides the order of the communication
- ▶ all processors wait for completion before proceeding
(source code – Fortran)

Communication Pattern 3



- blocking communication
 - odd/even MPI_SENDRECV
- alternating pairs communicate
 ⇒ needs only 2 steps

(source code – C)

Overlapping computation & communication

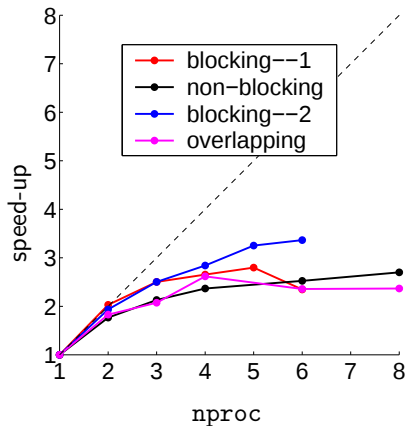
Symbolic code

```
(1) initialisation
for each iteration
  (2a) update local part of solution
  (2b) check for convergence
  (2c) exchange ghost-cell data
end
```

- ▶ overlap (2c) with (2a)
- ▶ non-blocking communication
- ▶ treat grid extrema after interior points,
check for termination of communication

(jacobi-step, communication)

Parallel Speed-Up (1)



2D Poisson problem:

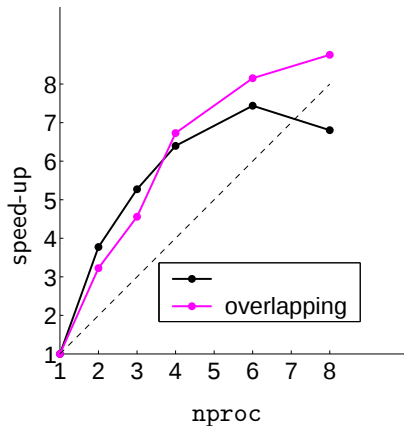
$n=m=250$

SGL Altix

shared memory

`fenix.ciemat.es`

Parallel Speed-Up (2)



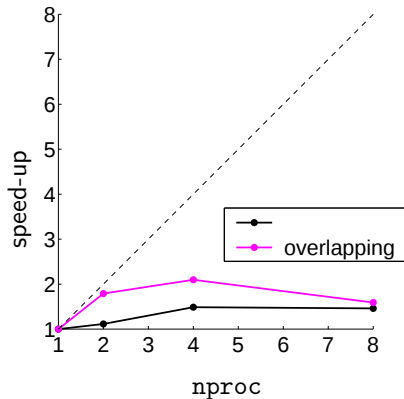
2D Poisson problem:

$n=m=500$

SGI Altix
shared memory
fenix.ciemat.es

super-linear

Parallel Speed-Up (3)



2D Poisson problem:

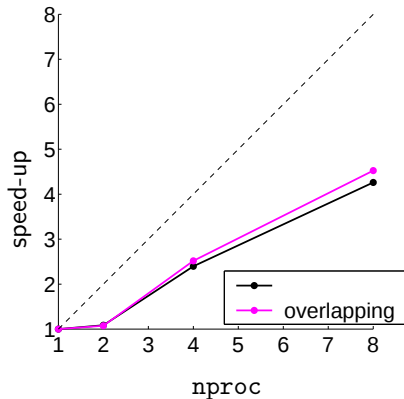
$n=m=250$

Xeon cluster

Gbit ethernet

`lince.ciemat.es`

Parallel Speed-Up (4)



2D Poisson problem:

$n=m=250$

Xeon cluster

InfiniBand

`lince.ciemat.es`

Parallel performance for 2D Poisson problem

Some observations

- ▶ speedup is problem-size dependent
often: larger size, more parallel work (Amdahl)
 - ▶ speedup is hardware dependent
CPU speed; network bandwidth/latency; memory bandwidth
 - ▶ super-linear speedup due to cache effects
 - ▶ design of communication patterns matters
 - ▶ overlapping communication **may** be beneficial
- ~> “optimal” solution differs from machine to machine

Acknowledgement

This lecture heavily draws from the following resource:

- ▶ <http://webct.ncsa.uiuc.edu:8900>
(Introduction to MPI)

MPI code analysis with MPE

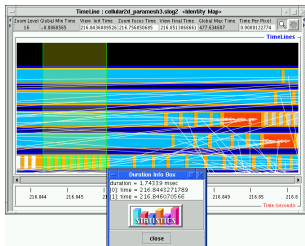
Multi Processing Environment

- ▶ library and utility programs distributed with MPICH
- ▶ content:
 - ▶ routines for creating logfiles
 - ▶ graphical visualization tools (jumpshot)
 - ▶ shared-display parallel X graphics library
 - ▶ miscellaneous other tools

MPI code analysis with MPE (2)

Automatic profiling

- compile code with linker flag `-mpilog`
- execution generates logfile `<execname>.clog`
- ⇒ then visually analyze the result with `jumpshot <execname>.clog`



Manual profiling example

```
PROGRAM HELLO  !compile with -lmpe
INCLUDE 'mpif.h'
INCLUDE 'mpef.h'
INTEGER myrank, size, ierr
C   Initialize MPI:
call MPI_INIT(ierr)
call MPI_COMM_RANK(MPI_COMM_WORLD,myrank,ierr)
call MPE_INIT_LOG()
call MPE_DESCRIBE_STATE(1,2,'hello','red:gray')
call MPE_LOG_EVENT(1,myrank,'write')
WRITE(*,*)"Hello_World!"
call MPE_LOG_EVENT(2,myrank,'written')
C   Terminate MPI, create log file "mpe_log.clog"
call MPE_FINISH_LOG("mpe_log")
call MPI_FINALIZE(ierr)
END
```

MPI code analysis with MPE

More information

- ▶ MPE library [documentation](#)
- ▶ Performance tool “jumpshot” [documentation](#)